

CS 521: Systems Programming

Inter-Process Communication

Lecture 15

Inter-Process Communication

- We previously discussed how the host OS tries its best to isolate processes
 - Processes should not be able to interfere with one another
 - To do *privileged* operations, we need to go through the kernel with **system calls**
- However, it's often useful to have processes communicate
 - Inter-Process Communication (IPC)
- IPC gives us safe, well-defined ways to communicate

Why IPC?

- Processes need to share data
- “Data” can mean a lot of things:
 - Plain text
 - An image, video, program
 - A message containing commands or other types of information
- Without a well-defined interface, getting processes to communicate descends into madness

An Example

1. You double-click a web link saved to your desktop
2. The OS determines which program is responsible for handling HTTP/S URIs
3. The program is launched if it isn't already running
4. The OS delivers a message to the program:
5. OPEN <https://google.com>

Types of IPC

- We will cover three types of IPC in this class (although there are many others):
 - Files
 - Signals
 - Pipes
- You might be surprised that files could be considered a form of IPC, but it's actually one of the easiest and simplest ways to communicate between processes!

Today's Schedule

- Files
- Signals
- Pipes

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Files

- Save a file to disk with one application, open it with another application
 - Needs a *file system* to make this happen
 - On our VMs, we're using `ext4` . A recent Mac might use `apfs` , and Windows `NTFS` (or maybe `XFAT` ...)
- What happens when two applications open the same file?
 - Coordinate via **file locks**
 - Can lock an entire file or only a portion

Opening a File

- We have used `fopen()` to open and read files
- Lower-level option: `open()`
 - This is a system call
 - (`fopen` from the C library calls `open` on Linux)
- `open` returns a *file descriptor* – an integer that represents the opened file
 - This decouples the file's absolute path in the file system (e.g., `/usr/bin/something`) from I/O operations

File Descriptors

- `stdin`, `stdout`, and `stderr` have file descriptors
- The file abstraction is used thoroughly in Unix systems (see `/dev` for devices)
- Once you've opened a file descriptor, you can read/write the contents of the file or even redirect the stream somewhere else

Redirecting Streams: dup2

- `dup2` allows us to redirect streams
 - `int dup2(int fildes, int fildes2);`
- Let's say we want to make our standard output stream go to a file we just opened
- We'll do:
 - `dup2(fd, STDOUT_FILENO);`
- This also deallocates (closes) the second fd
 - You won't see text printing directly to your terminal anymore

Example: Redirecting to a File

- Combine `open` and `dup2` :
 - `int output = open("output.txt",
O_CREAT | O_WRONLY | O_TRUNC, 0666);`
 - `dup2(output, STDOUT_FILENO);`
- This is exactly what our shell does when we use `<` and `>`
- `cat /etc/passwd > some_file`
 - Opens "some_file" and then redirects the output of the child process to that file instead!

Redirection Workflow

Let's say your shell encounters `>` in the command line...

1. Use `fork` to create a new process
2. Open the file that comes after the `>`
3. Redirect `stdout` to the file with `dup2`
4. Call `exec` to execute the program
 - This is part of the reason why `fork` and `exec` are split into two separate parts

Demo: io-redir.c

Today's Schedule

- Files
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- Pipes

Signals

- Signals are software-based *interrupts*
 - Basically a notification sent to the process
- The kernel uses signals to inform processes when events occur
- Handling a signal causes a jump in your program's logic to a *signal handler*
 - You can use a null signal handler to ignore particular signals

Demo: signal.c

Events

- What kind of events are reported via signals?
- It depends on the kernel
- To find out, use:

```
/bin/kill -l
```
- Wait, what?!
 - That's right: kill is used to send signals to processes
 - It doesn't necessarily 'kill' the process in doing so
 - But it can!

Terminating a Process

- You've already been using signals quite a bit (but maybe didn't realize)
 - Ever hit Ctrl+C to stop a running program?
 - it sends SIGINT to the process
- Each signal is prefixed with SIG
- Processes can choose how to deal with signals when they are received
 - Including ignoring them... usually

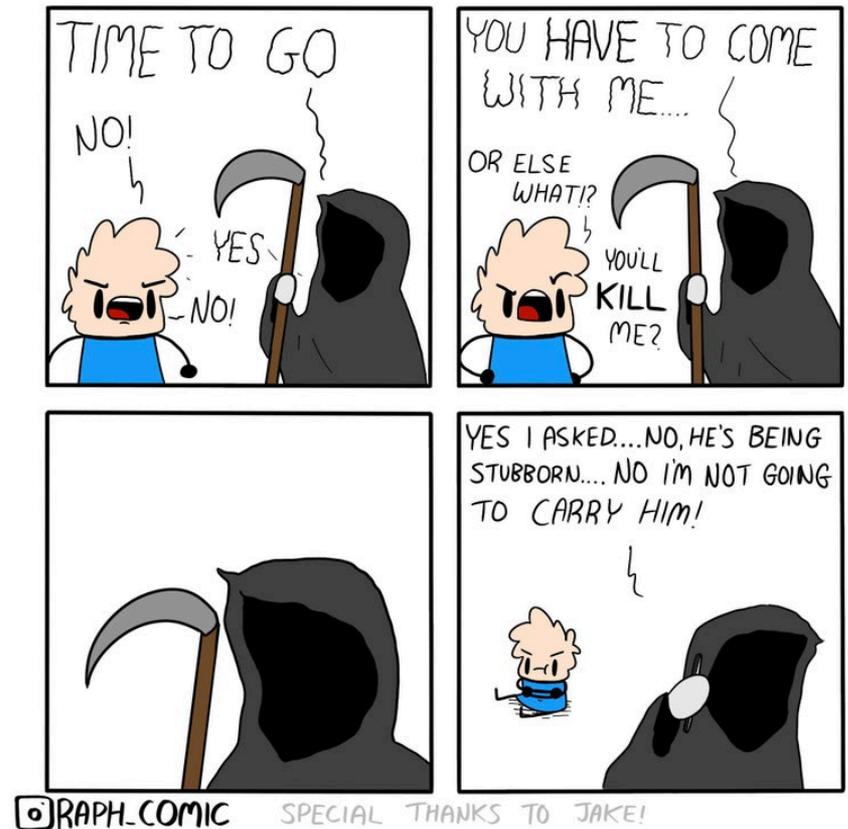
Demo: unkillable.c

Special Signals

- SIGSTOP and SIGKILL cannot be caught or ignored
- SIGSTOP – stops (pauses) the process: Ctrl+Z
 - SIGCONT tells a paused process to continue
- SIGKILL – terminates the process, no questions asked
 - You may have heard of `kill -9 <pid>`
 - 9 is SIGKILL

Using kill -9

- Occasionally a process will not respond to a SIGTERM, SIGINT, etc.
- This is the appropriate time to use SIGKILL



Signal Handling

- Set up a signal handler with `signal` :
`signal(SIGINT, sigint_handler);`
 - Will call `sigint_handler` every time a SIGINT is received
- Then implement the signal handling logic:
`void sigint_handler(int signo) { ... }`

OS Signal Transmission Process

1. First, a process initiates the signal
 - Terminal Emulator: user pressed Ctrl+C, so
 - I should send SIGTERM to the current process
2. The kernel receives the signal request
3. Permissions are verified
 - Can this user really send a signal to PID 3241?
4. The signal is delivered to the process

Reacting to a Signal

- If a process is busy doing something, it will be **interrupted** by the signal
- Jumps from the current instruction to the signal handler
 - (or performs the default operation if there is no signal handler)
- Jumps back to where it was when the handler logic completes

Segmentation Violation

- Our good friend, the segmentation violation (aka segfault) is also a signal
 - SIGSEGV
- Bus error: SIGBUS
- So if segfaults are getting you down, try blocking them!
 - What could go wrong?!

Sending a Signal

- Not all signals are sent via key combinations from the shell... We can send them programmatically or via the command line
- Let's send a SIGUSR1 signal to process 324:

```
kill -s SIGUSR1 324
```

- Simple as that!
- Or, in C:

```
int kill(pid_t pid, int signum);
```

Tracking Children

- `SIGCHLD` is sent to the parent of a child process when it exits, is interrupted, or resumes execution
- Useful in scenarios where the parent process needs to be notified about child events
 - or, in other words, when the parent is not already `wait()` ing on the child
 - Job list in the shell: when `SIGCHLD` is received, do a non-blocking `waitpid` to determine which process exited and remove it from the list (if backgrounded)

Today's Schedule

- Files
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- **Pipes**

Pipes [1/2]

- Pipes are a common way for programs to communicate on Unix systems
 - `cat /etc/something | sort | head -n5`
- Most useful for sharing unstructured data (such a text) between processes
- They work like how they sound: if you want to send data to another process, send it through the pipe

Pipes [2/2]

- Pipes are one of the fundamental forms of Unix IPC
- With pipes, we can “glue” several utilities together:
 - `grep neato file.txt | sort`
 - This will search for “neato” in file.txt and print each match
 - Next, these matches get sent over to the ‘sort’ utility
- Just like with I/O redirection, this is facilitated by `dup2`

In the Shell

- As we've seen, pipes are used frequently in the shell
- We can mix and match different utilities, and they all work well together
 - Awesome!
- Some genius must have designed all these programs to work this way, right?
 - Well, no. They all just read from `stdin` and then write to `stdout` (and `stderr`)
 - No coordination required between developers

Builtins vs. External Programs

- When you enter 'ls' in your shell, you're running a program
- This functionality is **NOT** built into your shell. Bash simply finds and runs the 'ls' program. That's it!
- There are some shell "commands" that actually aren't programs, called **built-ins**
 - `history`
 - `exit`
 - `cd` – why does this need to be a built-in?

Going to the Source

- I have posted a video from Bell Labs on the schedule that discusses several design aspects of Unix
- Discussion on pipes starts right around the 5 minute mark

The pipe function

- Now back to pipes: we can create them with the `pipe()` function
 - Returns a set of file descriptors: the **input** and **output** sides of the pipe
- Pipes aren't very useful if they aren't connected to anything, though
 - We can do this by `fork()` ing another process

Piping to Another Process

- After calling `fork()`, both processes have a copy of the pipe file descriptors
- Pipes only operate in one direction, though, so we need to close the appropriate ends of the pipe
 - You can think of a forked() pipe as one with four ends: two input and output ends each
 - We eliminate the ends we don't need to control the direction of data flow
 - Amazing ASCII art drawing: `>----<`

Controlling Flow

- To control data flow through the pipe, we close the ends we won't use
- For example:
 - Child process closes **FD 0** and reads from **FD 1**
 - Parent process closes **FD 1** and writes to **FD 0**

Async Process Creation

- You may be wondering: what good are pipes when we have to start all the cooperating processes?
- There's actually another option: **FIFOs**, aka **named pipes**
- Create with the `mkfifo` command, then open as you would a regular file descriptor

Redirecting Streams to a Pipe

- Let's say we want to make our standard output stream go through the pipe we just created
 - `int fd[2];`
`pipe(fd);`
- We'll do:
 - `dup2(fd[0], STDOUT_FILENO);`

Wrapping Up

- We've seen only a few possibilities for IPC!
- Another option: **sockets**
 - Communication... even over the network!
- Many Unix systems use *D-Bus* for more advanced IPC
- Windows has a similar concept: *Windows Messages*
 - Windows applications are **event based**
 - Almost everything that happens on Windows has an **event** associated with it (`WM_MOUSEMOVE` , changing resolution, etc.)

Fun: Undelivered Events

